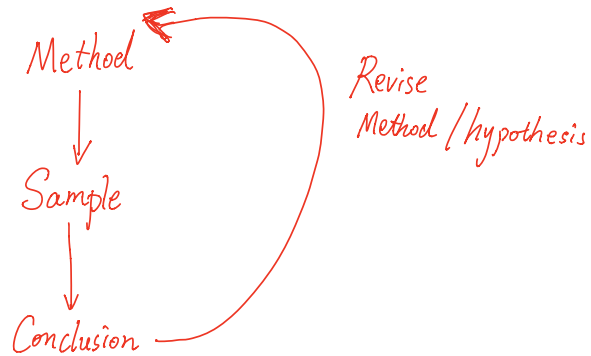


Lecture 23

Adaptive Data Analysis.



Logistics :

Project Presentation (May 3 & 5)

Schedule announced this week

20 mins.

No more HW.

Model for Adaptive Data Analysis

Statistical / Linear Queries

$$\phi: \mathcal{X} \rightarrow [0,1] \quad \text{"predicate"}$$

$$f_{\phi}(P) = \mathbb{E}_{x \sim P} [\phi(x)] \quad \text{"Population value"}$$

↑
population

$$f_{\phi}(D) = \mathbb{E}_{d_i \sim D} [\phi(d_i)] = \frac{1}{n} \sum_{i=1}^n [\phi(d_i)] \quad \text{"Empirical Average"}$$

$$D = (d_1, \dots, d_n) \in \mathcal{X}^n$$

Example: mean, correlation, variance, error/risk, gradients

Extension:

① Low-Sensitive Queries

$$\forall \text{ neighbors } D \& D', \quad |f_{\phi}(D) - f_{\phi}(D')| \leq \Delta$$

② Minimization Queries

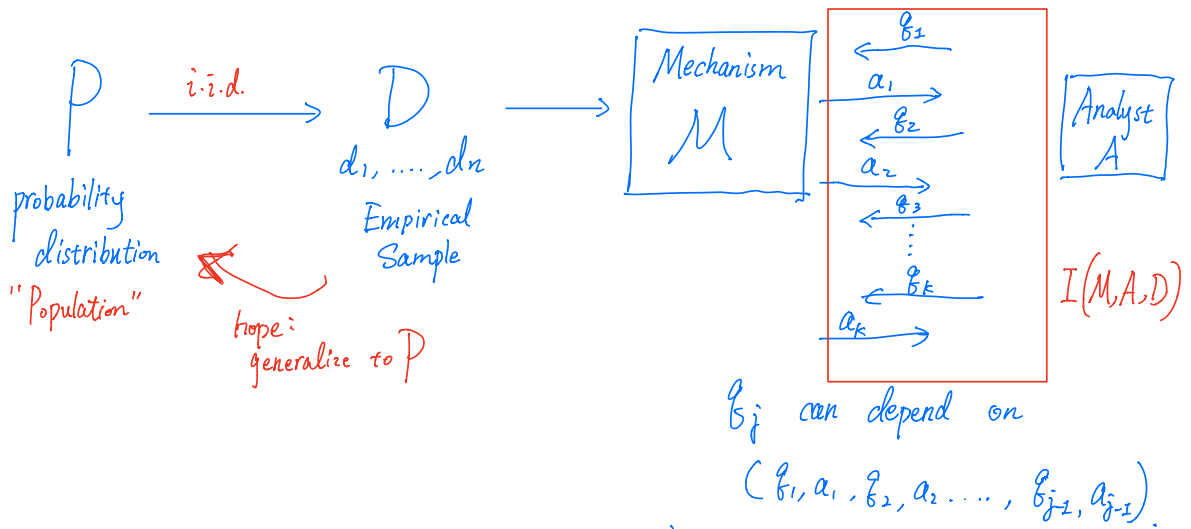
Query is given by some loss function $L: \mathcal{X}^n \times \Theta \rightarrow [0,1]$

↑ Dataset ↑ Parameter Space.

Answer $\theta \in \Theta$

$$\forall \text{ neighbors } D \& D', \quad \forall \theta \in \Theta, \quad |L(D, \theta) - L(D', \theta)| \leq \Delta.$$

Interaction of Adaptive Data Analysis.



Transcript $\Pi = (g_1, a_1, \dots, g_k, a_k) \leftarrow I(M, A, D)$

↑
Interaction

"Goal" = $\forall j,$

$$|a_j - g_j(P)| \leq \text{small.}$$

↑ Population Value.
Not just empirical averages

Avoid: queries g s.t.

$$|g(P) - g(D)| \geq \text{Large}$$

↑
 D is not representative.

DP \implies Generalization in ADA

(α, β) - sample accuracy

$$\mathbb{P}_{D \sim p^n, \pi} \left[\max_j |g_j(D) - a_j| \geq \alpha \right] \leq \beta$$

(α, β) - distributional accuracy

$$\mathbb{P}_{D \sim p^n, \pi} \left[\max_j |g_j(P) - a_j| \geq \alpha \right] \leq \beta$$

Sample Complexity / Accuracy.

Non-adaptive Queries.

- Take empirical averages: $a_j = f_j(D)$

$$\max_j |a_j - f_j(P)| \approx \sqrt{\frac{\log(k)}{n}}$$

Adaptive Queries

- Sample Splitting Method: D_1, \dots, D_k , $a_j = f(D_j)$

$$\max_j |a_j - f_j(P)| \leq \sqrt{\frac{k}{n}}$$

- Differential Privacy, Gaussian Mechanism

$$a_j = f_j(D) + N(0, \sigma^2).$$

$$\max_j |a_j - f_j(P)| \approx \underbrace{\frac{k^{\frac{1}{4}}}{\sqrt{n}}}_{\epsilon} + \underbrace{\sigma}_{\alpha} \quad O(\alpha + \epsilon).$$

Adding Noise
Reduces Error.

$$\underbrace{\frac{\sqrt{k}}{n\sigma}}_{\epsilon \text{ privacy/gen Bound}} + \underbrace{\sigma}_{\alpha \text{ sample Accuracy Bound}}$$

Transfer Theorem. (ϵ, δ) -version [JLNRSS20]

Suppose $I(M, A, D)$ is (α, β) -sample accurate \leftarrow
& (ϵ, δ) -DP. \leftarrow

Then for every $c, d > 0$, $I(M, A, D)$ is (α', β') \leftarrow
distributionally accurate, for

$$\alpha' = \alpha + \underbrace{(e^\epsilon - 1)}_{\approx \epsilon} + \underbrace{c}_{\alpha} + \underbrace{2d}_{\epsilon}, \quad \beta' = \frac{\beta}{c} + \frac{\delta}{d}$$

$$\alpha' = O(\alpha + \epsilon) \quad \beta' = \left(\frac{\beta}{\alpha} + \frac{\delta}{\epsilon} \right)$$

Simpler version $(\epsilon, 0)$ -DP

(α, β) -sample accuracy

$(\epsilon, 0)$ -DP.

$\Rightarrow (\alpha', \beta')$ -distributionally accurate.

$$\alpha' = \alpha + (e^\epsilon - 1) + \sqrt{\frac{2 \ln(1/\eta)}{n}}, \quad \beta' = \beta + \eta, \quad \text{for all } \eta > 0.$$

$$\approx \alpha + \epsilon + O\left(\frac{1}{\sqrt{n}}\right)$$

\leftarrow Sampling error bound

Proof Sketch.

$$D \stackrel{\text{i.i.d.}}{\leftarrow} \mathcal{P}^n$$

$$\text{Transcript: } \pi = (g_1, a_1, \dots, g_k, a_k) \leftarrow I(M, A, D)$$

$$Q_\pi = \mathcal{P}^n \mid \pi$$

"posterior distribution over D conditioned on π "

- Suppose M is (α, β) -sample accurate.

$$\mathbb{P} \left[\max_j |a_j - g_j(D)| \geq \alpha \right] \leq \beta$$

Event E about (D, π)

Lemma (Bayesian Resampling)

$$\mathbb{P}_{\substack{D \sim \mathcal{P}^n \\ \pi \leftarrow I(M, A, D)}} \left[(D, \pi) \in E \right] = \mathbb{P}_{\substack{D \sim \mathcal{P}^n \\ \pi \sim I(M, A, D) \\ D' \sim Q_\pi}} \left[(D', \pi) \in E \right]$$

- Generic: Nothing to do w/ DP.

- Sample accuracy w.r.t. D

\Rightarrow Sample accuracy w.r.t. $D' \sim Q_\pi$

- $f(D) \approx f(Q_\pi) = \mathbb{E}_{D' \sim Q_\pi} [f(D')]$

- Goal: $f(D) \approx f(P)$
 $\alpha \approx f(P)$

Missing Step: $f(Q_\pi) \approx f(P)$

Do you want to see the proof?

Bayesian Resampling
Lemma

$$\text{Proof. } \mathbb{P} \left[(D', \pi) \in E \right]$$

$D \sim p^n$
 $\pi \leftarrow I(M, A, D)$
 $D' \sim Q_\pi$

$$= \sum_{D=x} \sum_{\pi=\pi} \sum_{D'=x'} \mathbb{1}[(x', \pi) \in E] \cdot \mathbb{P}[D=x] \cdot \mathbb{P}[\pi|x] \cdot \mathbb{P}[x'|\pi]$$

$$= \sum_{D=x} \sum_{\pi=\pi} \sum_{D'=x'} \mathbb{1}[(x', \pi) \in E] \cdot \mathbb{P}[D=x, \pi=\pi] \mathbb{P}[D'=x' | \pi=\pi]$$

$D \sim p^n$

$$= \sum_{\pi=\pi} \sum_{D'=x'} \mathbb{1}[(x', \pi) \in E] \cdot \mathbb{P}[\pi=\pi] \cdot \mathbb{P}[D'=x' | \pi=\pi]$$

$\pi \leftarrow I(M, A, D)$
 $D \sim p^n$
 $\pi \leftarrow I$

$$= \sum_{\pi=\pi} \sum_{D'=x'} \mathbb{1}[(x', \pi) \in E] \cdot \mathbb{P}[D'=x'] \cdot \mathbb{P}[\pi=\pi | D'=x']$$

$D \sim p^n$
 $\pi \leftarrow I$

$$= \mathbb{P}[(D, \pi) \in E]$$

$D \sim p^n$
 $\pi \leftarrow I(M, A, D)$

Differential Privacy $\Rightarrow f(Q_\pi) \approx f(P)$

Proof (Sketch).

$$\begin{aligned}
 f(Q_\pi) &= \mathbb{E}_{D' \sim Q_\pi} [f(D')] \\
 &= \mathbb{E}_{\substack{D' \sim Q_\pi \\ i \leftarrow \text{unif}\{1, \dots, n\}}} [f(d'_i)] = \mathbb{E}_{\substack{D \sim P^n \\ i \leftarrow \text{unif}\{1, \dots, n\}}} [f(d_i)] \\
 &= \int_{x \in X} f(x) \cdot \mathbb{P}_{\substack{D \sim P^n \\ i \leftarrow \text{unif}\{1, \dots, n\}}} [d_i = x \mid \pi] \quad \leftarrow \begin{array}{l} \text{Conditioning} \\ \text{Bayes} \\ \text{Rule} \end{array} \\
 &= \int_{x \in X} f(x) \cdot \underbrace{\mathbb{P}_{\substack{D \sim P^n \\ i \leftarrow \text{unif}\{1, \dots, n\}}} [\pi \mid d_i = x]}_{\substack{D \sim P^n \\ i \leftarrow \{1, n\}}} \cdot \mathbb{P}_{\substack{D \sim P^n \\ i \leftarrow \{1, n\}}} [d_i = x] \\
 &\leq \int_{x \in X} f(x) \frac{e^\epsilon \mathbb{P}[\pi] \cdot \mathbb{P}[d_i = x]}{\mathbb{P}[\pi]} \quad \leftarrow \begin{array}{l} \text{Approximate} \\ \text{Cancellation?} \\ \pi \leftarrow I(M, A, D) \\ \uparrow \\ (\epsilon, \delta)\text{-DP} \end{array} \\
 &= e^\epsilon \int_{x \in X} f(x) \cdot \mathbb{P}_{\substack{D \sim P^n \\ i \sim \{1, \dots, n\}}} [d_i = x] = f(P).
 \end{aligned}$$

$$f(Q_\pi) \leq e^\epsilon f(P)$$

$$f(Q_\pi) \geq e^{-\epsilon} f(P).$$

Putting Together

- $f(Q_n) \in [e^{-\epsilon} f(P), e^{\epsilon} f(P)]$

Due to (ϵ, δ) -DP.

- " $f(D) \approx f(Q_n)$ "

$$\mathbb{P}_{\substack{D \leftarrow P^n \\ z \leftarrow I}} [|f(Q_n) - f(D)| \geq C_\eta] = \mathbb{P}_{\substack{D \leftarrow P^n \\ z \leftarrow I \\ D' \leftarrow Q_n}} [|f(Q_n) - f(D')| \geq C_\eta] \leq \eta$$

$$C_\eta = \sqrt{\frac{2 \ln(1/\eta)}{n}}$$

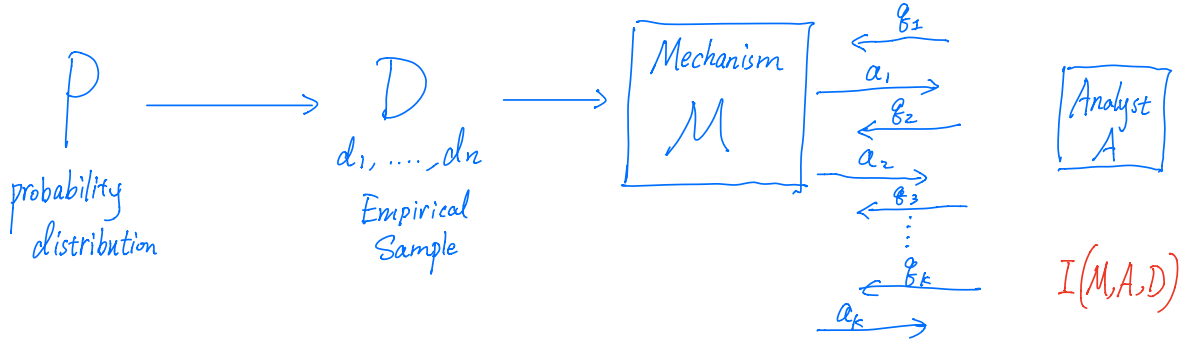
Azuma
Ineq.

- Sample accuracy, w.p. $1 - \beta$.

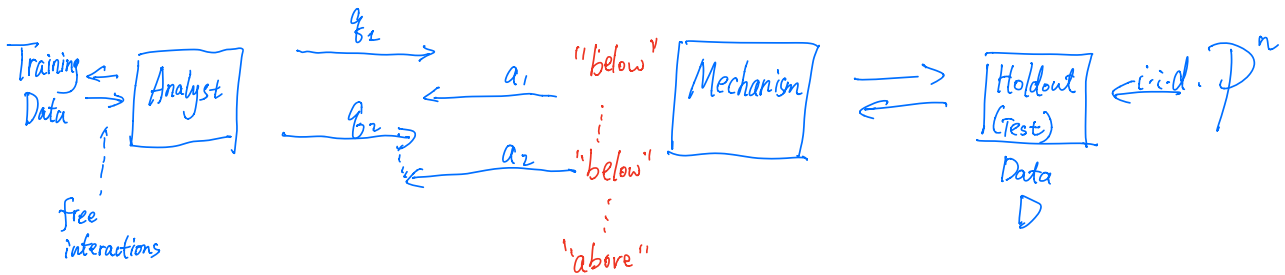
$$|a - f(D)| \leq \alpha \quad \text{for all queries. } f = f_j.$$

Add 3 sources of error \square

Mental Model of ADA.



b_j can depend on $(b_1, a_1, b_2, a_2, \dots, b_{j-1}, a_{j-1})$.



Example:
$$g_1(D) = \frac{1}{n} \sum_{i=1}^n \mathbb{I}[h(x_i) = y_i]$$

\uparrow
~~error~~
 accuracy

Reusable Holdout.

Sparse Vector or Thresholdout

Dwork et al.
Science '2016

Algorithm 2: SparseVector($s, T, \Delta, \epsilon, q_1, q_2, \dots$):

Input: q_1, q_2, \dots is a stream of Δ -sensitive queries

- 1 **AllDone** \leftarrow **FALSE**; T : threshold
- 2 $\tilde{T} = T + Z_0$ where $Z_0 \sim \text{Lap}(2\Delta/\epsilon)$; Stop when $q(s)$ is above T .
"L(0)"
- 3 **while not AllDone do**
- 4 Accept the next query q_i ;
- 5 $a_i \leftarrow q_i(s)$;
- 6 $\tilde{a}_i \leftarrow a_i + Z_i$ where $Z_i \sim \text{Lap}(4\Delta/\epsilon)$;
- 7 **if** $\tilde{a}_i < \tilde{T}$ **then**
- 8 **return** $b_j = \perp$;
- 9 **else**
- 10 **return** $b_j = \top$;
- 11 **AllDone** \leftarrow **TRUE**;

- Only Noise T once.
- Could release " \perp " many times.